

4.4 Counting Lemma

The *homomorphism density* $t(F, G)$ of a simple graph F into a simple graph G , is defined as the probability that a uniform random map of $V(F)$ into $V(G)$ preserves the edges. In formula

$$t(F, G) = \frac{\text{hom}(F, G)}{|V(G)|^{|V(F)|}},$$

where $\text{hom}(F, G)$ is the number of homomorphisms (edge-preserving maps) $V(F) \rightarrow V(G)$. As an example, $t(K_2, G) = 2|E|/|V|^2$.

Lemma 4.4.1 *For any three simple graphs F , G_1 and G_2 ,*

$$|t(F, G_1) - t(F, G_2)| \leq |E(F)| \delta_{\square}(G_1, G_2).$$

Proof. Let ϕ be a uniform random map $V(F) \rightarrow V(G)$. Let $E(F) = E_1 \cup E_2$ be a partition of $E(F)$ into two (possibly empty) sets. For $X \subseteq E$, consider the probability

$$p(X) = \mathbf{P}(\phi(X) \subseteq E(G_1), \phi(E \setminus X) \subseteq E(G_2)).$$

Clearly $p(E) = t(F, G_1)$ and $p(\emptyset) = t(F, G_2)$. It suffices to prove that if $X' = X \cup \{ij\}$ ($ij \in E \setminus X$), then $|p(X') - p(X)| \leq \delta_{\square}(G_1, G_2)$.

Let $V' = V \setminus \{i, j\}$ and fix the map $\phi' = \phi|_{V'}$. Node i can be mapped onto a set $S \subseteq V$ so that the edges between i and V' are mapped properly (this is the same for X and X'), and similarly j can be mapped onto a set $T \subseteq V$. The edge ij must be mapped into $E(G_1)$ to contribute to $p(X')$, but into $E(G_2)$ to contribute to $p(X)$. So

$$\mathbf{P}(\phi(X) \subseteq E(G_1), \phi(E \setminus X) \subseteq E(G_2)) | \phi' = \frac{e_{G_2}(S, T)}{n^2},$$

and

$$\mathbf{P}(\phi(X') \subseteq E(G_1), \phi(E \setminus X') \subseteq E(G_2)) | \phi' = \frac{e_{G_1}(S, T)}{n^2}.$$

It follows that

$$|\mathbf{P}(\phi(X) \subseteq E(G_1), \phi(E \setminus X) \subseteq E(G_2)) | \phi' - \mathbf{P}(\phi(X') \subseteq E(G_1), \phi(E \setminus X') \subseteq E(G_2)) | \phi'| = \left| \frac{e_{G_2}(S, T)}{n^2} - \frac{e_{G_1}(S, T)}{n^2} \right| \leq$$

Since this holds for every ϕ' , the lemma follows. \square